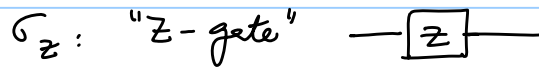


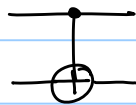
Note: In a quantum circuit, Paulis are represented as:



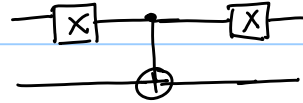
* Measurement in the quantum circuit model:-

(i) Deferred Measurement:- Measurements can always be moved to the end of a circuit by use of controlled operations.

Using CNOT:



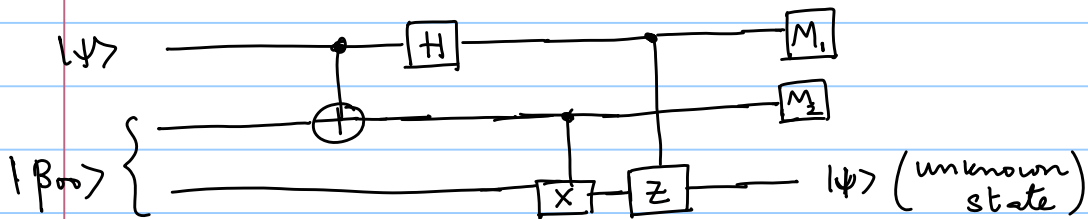
(or)



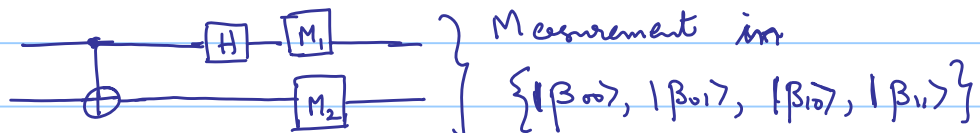
Flips target if control is set to '1'

Flips target if control is set to zero!

Eg. Teleportation:-



(ii) Measuring in the Bell basis:-



* Universal Quantum Gates

Defn: A set of gates $G = \{g_1, g_2, \dots, g_n\}$ is universal for quantum computation if any unitary operation can be approximated to arbitrary accuracy by a quantum circuit using only those gates.

One set of universal gates: $\{H, S, CNOT, G\pi/8\}$.

* Two-step proof:-

(1) Arbitrary U can be exactly realized as a product of single-qubit unitaries and CNOT

(2) Any single-qubit unitary can be approximated to arbitrary accuracy using $H, S, G\pi/8$.

(1) Exact Universality:-

Step (1a): Arbitrary $d \times d$ unitary \rightarrow 2-level unitary
(acts non-trivially on some 2-dim subspace of d -dim space)

Eg. Consider the 3-qubit gate:-

$$U = \begin{pmatrix} a & 0 & 0 & 0 & 0 & 0 & 0 & c \\ 0 & 1 & & & & & & 0 \\ 0 & & 1 & & & & & 0 \\ 0 & & & 1 & & & & 0 \\ 0 & & & & 1 & & & 0 \\ 0 & 0 & & & & 1 & & 0 \\ 0 & & & & & & 1 & 0 \\ b & 0 & 0 & 0 & 0 & 0 & 0 & d \end{pmatrix}$$

Acts non-trivially only on span of $\{|000\rangle, |111\rangle\}$

We show that:- Say U is a $d \times d$ unitary,

\exists 2-level unitaries $U_1, U_2, U_3 \dots$ etc

such that $U_k U_{k-1} U_{k-2} \dots U_3 U_2 U_1 U = I$

$$\therefore U = U_1^\dagger U_2^\dagger U_3^\dagger \dots U_k^\dagger$$

Example:- $U = \begin{pmatrix} a & b & c \\ d & e & f \\ g & h & j \end{pmatrix} \quad (3 \times 3)$

(i) If $d=0$, let $U_1 = \begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix}$

otherwise, $U_1 = \begin{pmatrix} \frac{a^*}{\sqrt{|a|^2 + |d|^2}} & \frac{d^*}{\sqrt{|a|^2 + |d|^2}} & 0 \\ \frac{d}{\sqrt{|a|^2 + |d|^2}} & \frac{-a}{\sqrt{|a|^2 + |d|^2}} & 0 \\ 0 & 0 & 1 \end{pmatrix}$

$\left(\begin{pmatrix} a & b & c \\ d & e & f \\ g & h & j \end{pmatrix} \right)$

$$U_1 U = \begin{pmatrix} a' & b' & c' \\ 0 & e' & f' \\ g & h & j \end{pmatrix}$$

(ii) Why if $g=0$, $U_2 = \begin{pmatrix} a'^* & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix}$, otherwise,

$$U_2 = \begin{pmatrix} \frac{a'^*}{\sqrt{|a'|^2 + |g|^2}} & 0 & \frac{g^*}{\sqrt{|a'|^2 + |g|^2}} \\ 0 & 1 & 0 \\ \frac{a}{\sqrt{|a'|^2 + |g|^2}} & 0 & \frac{-g}{\sqrt{|a'|^2 + |g|^2}} \end{pmatrix}$$

$$\therefore U_2 U_1 U = \begin{pmatrix} 1 & b'' & c'' \\ 0 & e'' & f'' \\ 0 & h'' & j'' \end{pmatrix} = V$$

Since V is unitary, $b'' = c'' = 0$

$$\therefore U_2 U_1 U = \begin{pmatrix} 1 & 0 & 0 \\ 0 & e'' & f'' \\ 0 & h'' & j'' \end{pmatrix} = U_3^\dagger$$

$$\therefore U = U_1^\dagger U_2^\dagger U_3^\dagger$$

* Generally, for a $d \times d$ unitary U_d ,

using $(d-1)$ two-level matrices $\rightarrow \begin{pmatrix} 1 & 0 \\ 0 & \boxed{U_{d-1}} \end{pmatrix}$

$U_{d-1} \rightarrow (d-2)$ two-level matrices $\rightarrow \begin{pmatrix} 1 & 0 \\ 0 & \boxed{U_{d-2}} \end{pmatrix}$

etc.

$$\therefore U_d = V_1 V_2 \dots V_k \text{ (two-level unitaries)}$$

$$\text{with } k \leq (d-1) + (d-2) + \dots + 1$$

$$= \frac{d(d-1)}{2}$$

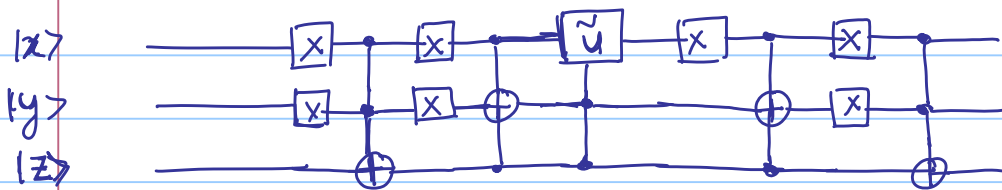
Step (1b):- Any two-level unitary can be realized exactly using single-qubit unitaries and CNOT.

Exact universality: $\mathcal{G} = \left\{ \begin{array}{l} \text{single-qubit} \\ \text{unitary} \\ \text{SU}(2) \end{array}, \text{CNOT} \right\}$

Example:

$$U = \begin{pmatrix} a & & & b \\ & 1 & & \\ & & 1 & \\ c & & & d \end{pmatrix}_{8 \times 8}$$

$$\tilde{U} = \begin{pmatrix} a & b \\ c & d \end{pmatrix} \text{ acting on } \{|0000\rangle, |1111\rangle\}$$



$$000 \rightarrow 001 \rightarrow 011$$

* Step (2): - Approximate universality:

• Quantifying error:

When V is implemented instead of U : $e(U, V)$

$$e(U, V) = \max_{|\psi\rangle} \|(U - V)|\psi\rangle\|,$$

- Recall, $\| |\phi\rangle \| = \sqrt{\langle \phi | \phi \rangle}$.

- If $e(U, V)$ is small, any measurement on $V|\psi\rangle$ gives approximately the same statistics as a measurement on $U|\psi\rangle$.

Let M_i be an arbitrary element of a POVM.

$p_U(i) \equiv$ prob. of outcome 'i' for the state $U|\psi\rangle$

$p_V(i) \equiv$ " " for state $V|\psi\rangle$.

then, $|p_U(i) - p_V(i)| \leq 2e(U, V)$.

Proof: $|\rho_U(i) - \rho_V(i)| = |\langle \psi | U^\dagger M_i U | \psi \rangle - \langle \psi | V^\dagger M_i V | \psi \rangle|$

Let $|\xi\rangle \equiv (U - V)|\psi\rangle$.

Then, $|\rho_U(i) - \rho_V(i)| = |\langle \psi | U^\dagger M_i |\xi\rangle + \langle \xi | M_i V | \psi \rangle|$

(triangle inequality) $\leq |\langle \psi | U^\dagger M_i |\xi\rangle| + |\langle \xi | M_i V | \psi \rangle|$

(Cauchy-Schwarz) $\leq \|U|\psi\rangle\| \|\xi\rangle\| + \|\xi\rangle\| \|V|\psi\rangle\|$

$= 2 \|(U - V)|\psi\rangle\| = 2 e(U, V)$.

- Further, the errors add up linearly for a sequence of unitaries.